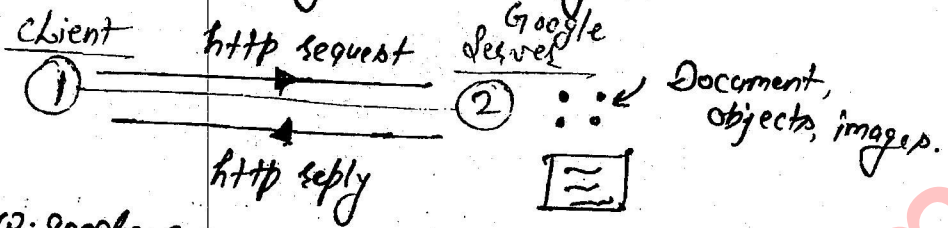


* Application Layer :-

⇒ "http Protocol" :- (hypertext transfer protocol)



http://www.google.com

url: uniform resource locator

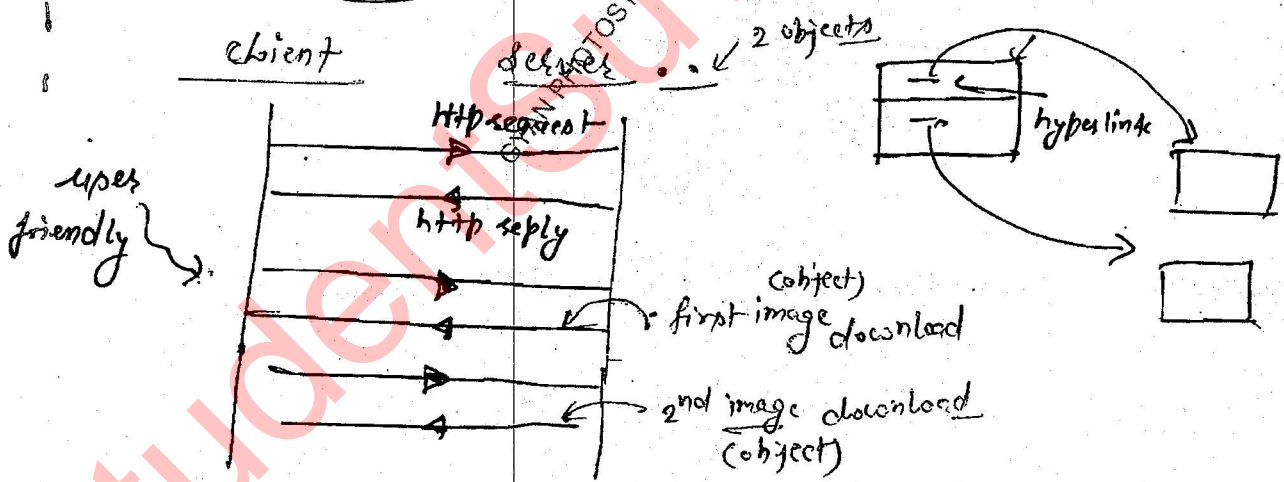
→ T.L. (Transport Layer)
S.Port 65000, D.Port 80 ← http

* http is a synchronous Protocol because both the clocks of client & server are synchronized.

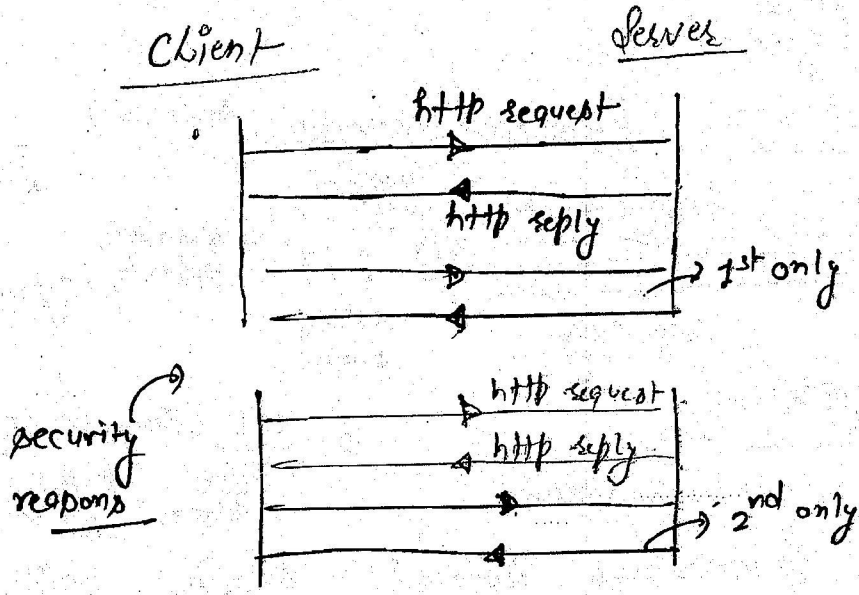
* http Connections :-

- ↳ (i) Persistent http Connection
- (ii) Non-persistent http connection

:- Persistent http :-

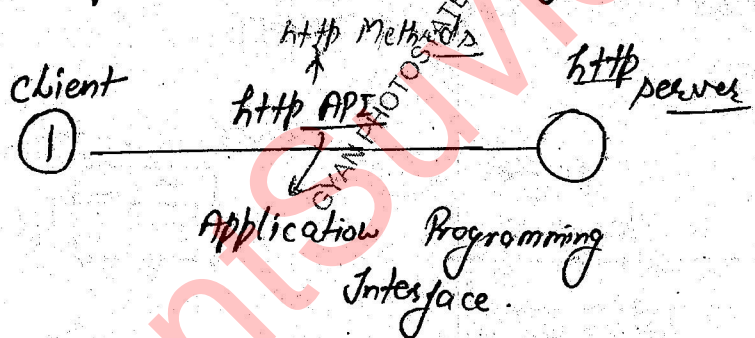


* In case of persistent http connection will be there only until all objects are accessed.
 * this is for user friendly requirement.



* In Non-persistent http, a separate connection is established for each every individual data transfer. This is provided for security.

http API :-



Methods:- get, post, put, head, trace, connect

* connect: http $\xrightarrow{\text{connect}}$ https

- * get method is used to retrieve the document
- * Post method is used to place the updated document in the server
- * Put method is used to modify the content of the document.
- * head method is used to get the information about the document

* Connect method is used to transmit the data in secure channel and that to in encrypted form.

* http is stateless protocol.

* By default IPv6 is stateless but we can make it stateful with the help of DHCPv6 server.

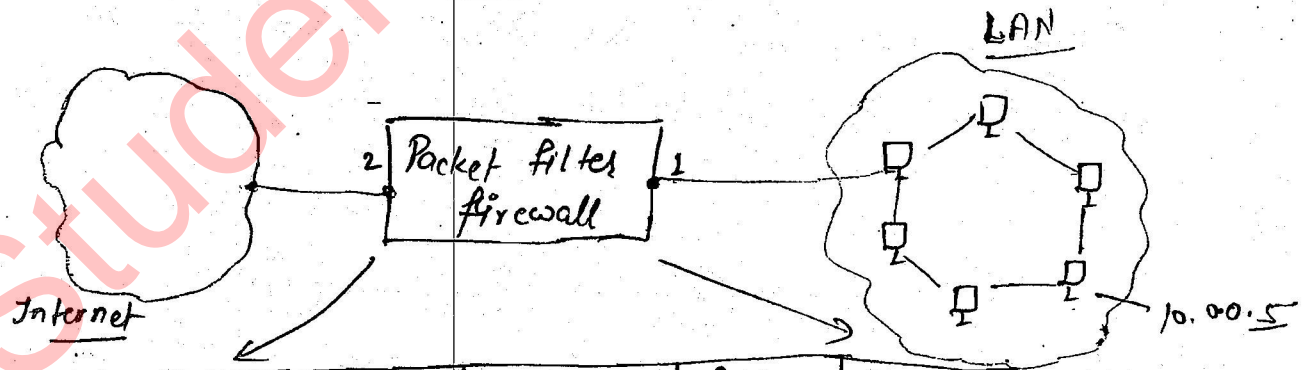
Note:- http is a stateless protocol because client will not store any information about the server once the transaction has been done.

* Cookie is a piece of code that is transmitted from the server to the client browser.

* The advantage of cookies is:
 (i) faster response
 (ii) authentication.

* "Firewall" :-

:- Packet filter firewall :-



Interface	Source IP	Source Port	Dest IP	Dest. Port
2	172.12.0.0	*	*	*
2	*	*	10.0.0.5	*
2	*	*	*	23
1	*	*	*	80

→ Telnet Port

* Packet filter firewall is a firewall which checks or forwards the data by observing the transport layer and Network layer 'headers' of the content by comparing with firewall table.

* There is no ideal firewall, every firewall will work according to its design.

* Packets coming from a particular source id i.e. 170.12.0.0 are blocked.

* Packets destined to 10.0.0.5 are blocked because this computer is used for internal LAN communication only.

* Packets destined to 'port-23' are blocked i.e. "Telnet" service is blocked. (No outside person can communicate directly to the internal LAN).

* When a malicious software or a virus is placed in the application data then the packet filter firewall can not detect it.

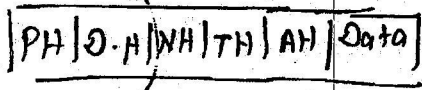
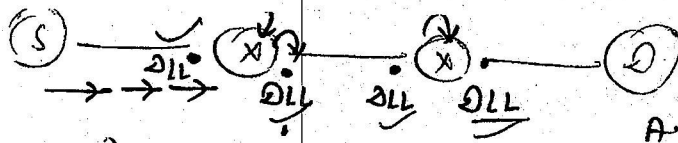
* Encryption and decryption is done in session layer ~~where~~ in OSI Model where as it is done in Application layer in TCP/IP Model.

①



Calculate how many times data link layer, N.L. are visited in this path from source to destination.
(TCP/IP)

A.L.
T.L
N.L ✓
DLL ✓
P.L

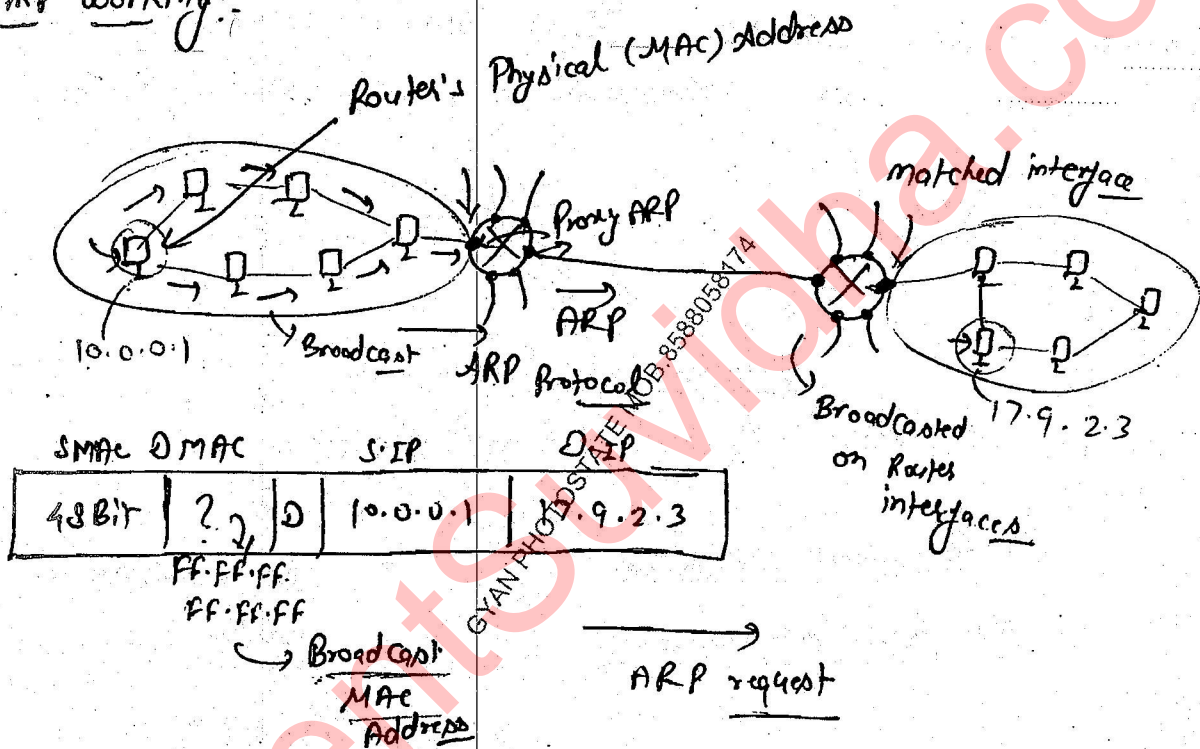


A.L
T.L
N.L ✓
DLL ✓
P.L

N.L = 4 times

D.L.L = 6 times

ARP working:



* FTP :- (file transfer Protocol)

has No flow control at A.L.

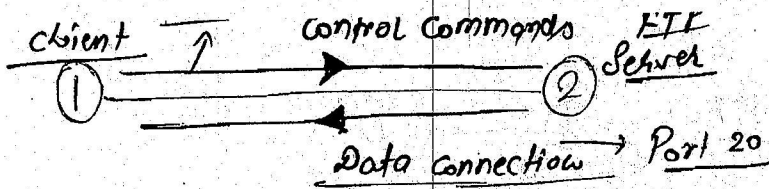
Reliable →
authorized user
TCP op T.L.

- (i) Internet
- (ii) purchasing reliable service

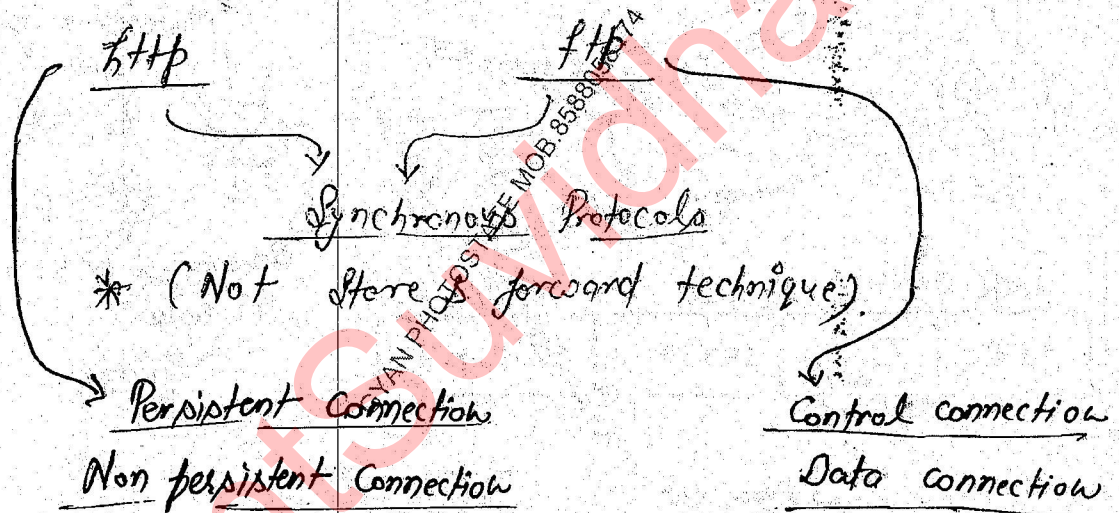
has flow control at A.L.

unreliable data
Anonymous user
UDP op T.L.

- (i) Internet Service
- (ii) Antivirus Software



- * FTP will send control commands via a control connection on Port - 21.
- * When the file is about to download, separate data connection will be established on Port - 20.
- * Once the file is completely downloaded, data connection is closed but control connection will be there, to download some other files.



FTP :-

- i) downloading a large file
- ii) Control connection
Data connection
- iii) Ports 20, 21 are used

TELNET :-

- i) chat operations (Exchange of words)
- ii) Common connection.
- iii) Port 23

TEL as T.L.

SMTP :- (simple mail transfer protocol)

(1) text-based Protocol

→

Internet → MIME Extension

↳ multimedia internet mail Extension

(SMTP + MIME) = all type of data

* Port - 25

Base 64 Encoding

0-25 ⇒ (0-25)

A-Z ⇒ (26-51)

0-9 ⇒ (52-61)

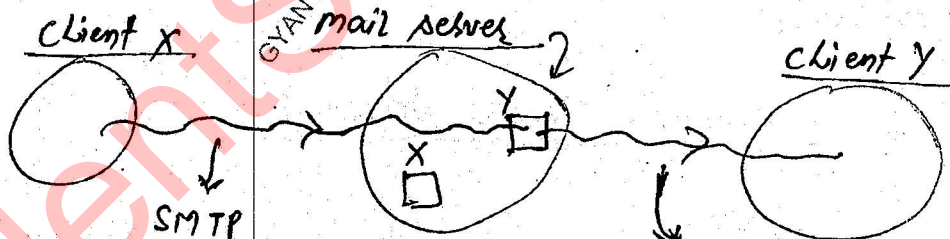
+ ⇒ (62, 63)

11 000 000 100 0000

↓ 48 10 0
ω K a

* TCP as T.L.

*



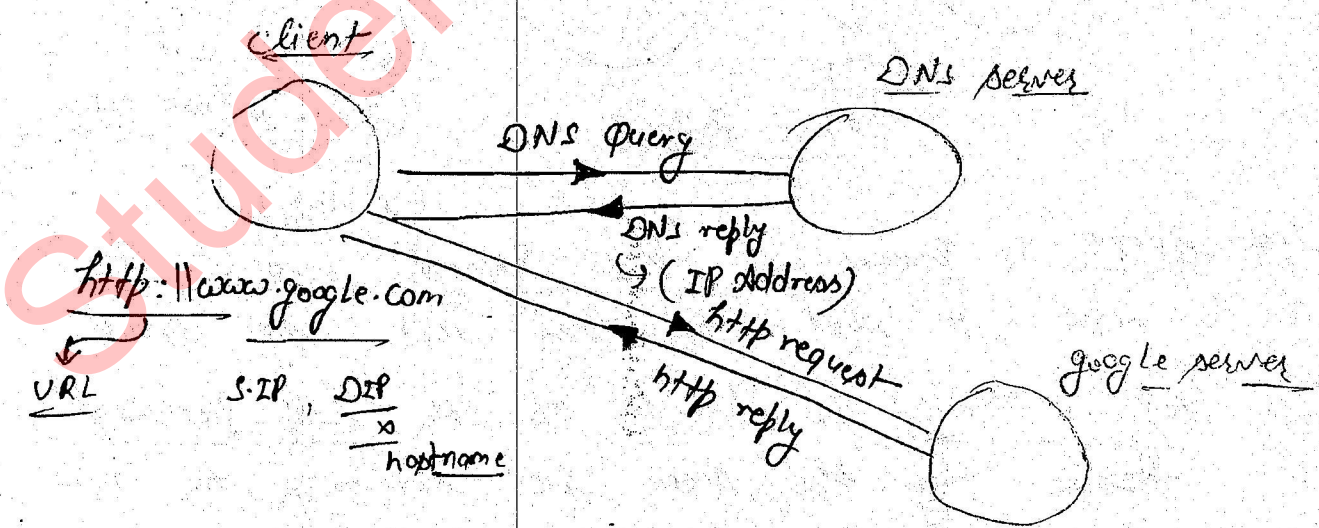
* SMTP is known as 'Push Protocol' because it is used to send the mail into mail server.

* POP₃ or IMAP₄ are known as the 'pull protocols' because they are used for retrieving the mails from mail server.

IMAP₄ :- (Internet message Access Protocol)

- * SMTP combined with POP3 is known as Client to Client Protocol with the mediation done by mail server.
- * SMTP with POP3 is applying store & forward technique.
- * SMTP combined with POP3 is a 'asynchronous Protocol' because their clocks need not to be synchronized.
- * In case of SMTP multiple paths are possible when and TCP is used as T.L. because it is a session.
- * In POP3 all mails are equal whereas in IMAP4 mails are kept in hierarchy.
- * IMAP4 is more secure than POP3, because it will scan for viruses before the file gets downloaded.

DNS : (Domain Name Space)



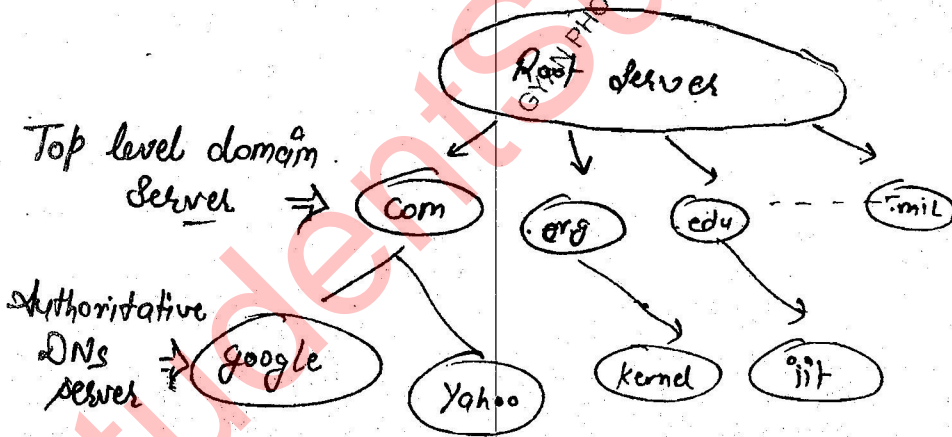
* The purpose of DNS is to find the destination IP address for the required hostname.
 Or it is used for mapping hostnames to IP-Addresses.

Design of DNS Server :-

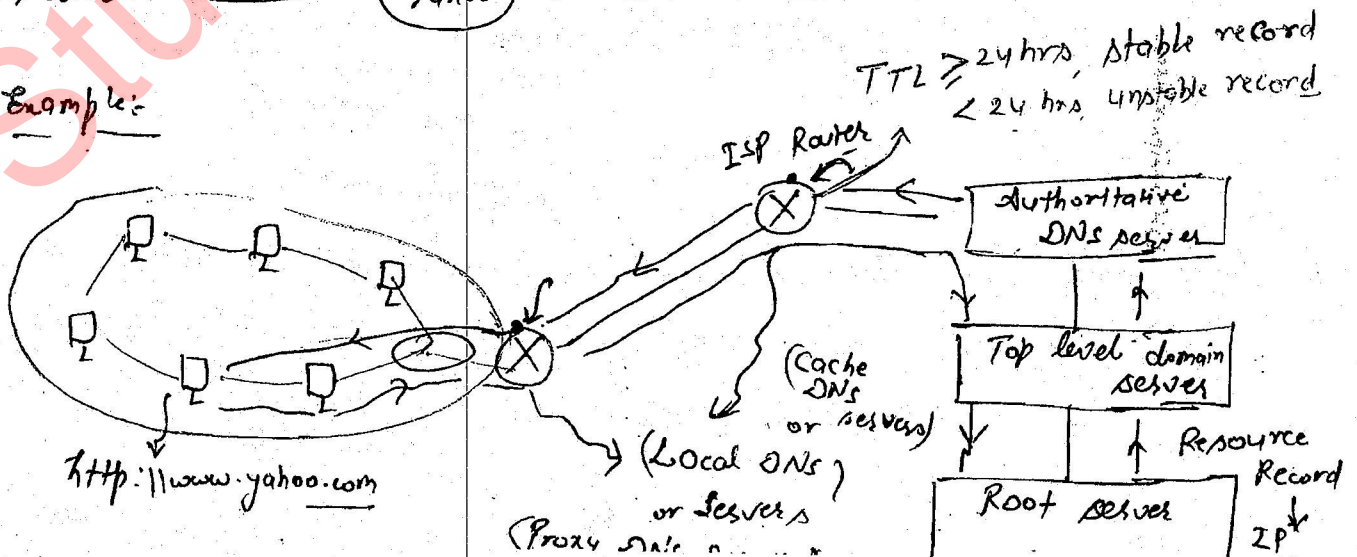
- (i) DNS servers placed in Hierarchy → Searching time will be less
- (ii) DNS Servers placed in different geographical locations → P.T. will be less

DNS servers

- Root Server
- Top level domain server
- Authoritative server
- Local DNS Server



Example:-



* DNS uses UDP as a transport layer protocol.

* The DNS ^{Reply} Record, it is stored for > 24 hours, it is treated as the stable record.

Que.:-



DNS

query size < 512 Bytes

→ UDP as T.L.

query size ≥ 512 Bytes

→ TCP as T.L.

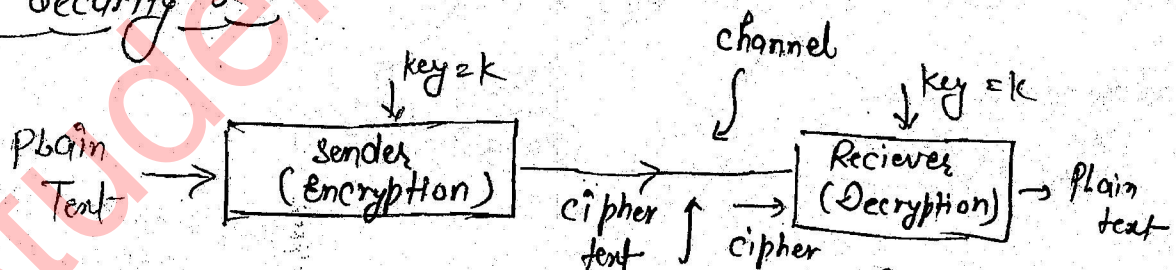
* In the DNS Server, there are two zones :-
'forward lookup zone', 'Reverse Lookup zone'.

* 'forward lookup zone' is used for mapping Host Names to IP-Addresses.

* Reverse lookup zone is used for mapping IP Addresses to Host Name.

Dated :- 21/06/2017

Basic Security :-



⇒ "Cryptography is the science or art of converting one form of data into other form, for providing security to the data is known as cryptography."

* Converting Plain text to cipher text is known as Encryption.

* Converting the ciphertext into plain text is known as Decryption.

* Steganography is the science of hiding the data behind an image or a video.

* The key is transmitted on the channel and later it is used for Encryption & decryption, it is treated as Public key.

* If the key is kept a secret and later it is used for Encryption and decryption, it is known as Private key.

Cryptography

Symmetric key
Cryptography
Ex: Diffie Hellman
Key Exchange

Asymmetric key
Cryptography
Ex: 'RSA Algorithm'

* In symmetric key cryptography same key is used for both Encryption and decryption.

* In asymmetric key cryptography different keys are used for encryption & decryption.

Challenges of Cryptography :-

(i) Authentication \rightarrow authentication of user
 \rightarrow authentication of data.

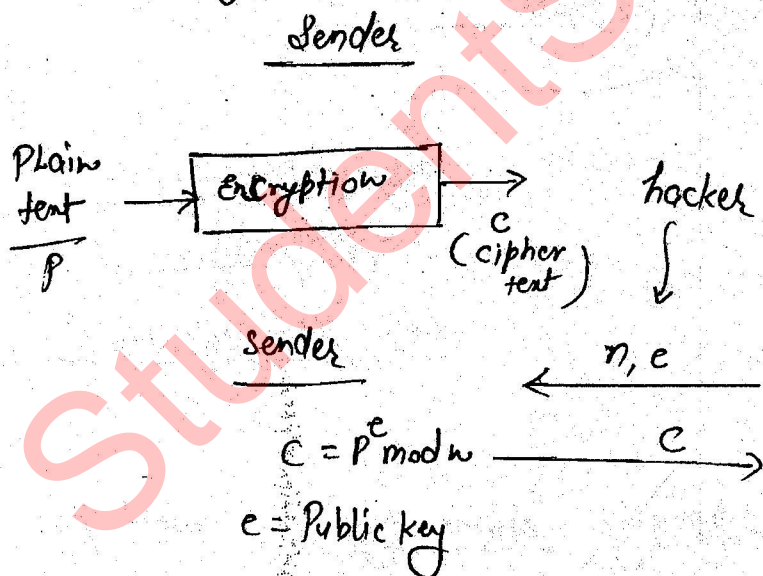
(ii) Confidentiality

- * Proving user's identity is known as 'authentication'.
- * Providing secrecy to the data is known as 'Confidentiality'.

Key features of cryptography :-

- ① Prime Numbers
- ② Challenges (Random number)
- ③ key values (Public key & Private key).
- ④ Timestamp.

RSA Algorithm :-



Receiver

P, q are Prime Numbers

$$n = P * q$$

$$\phi(n) = (P-1)(q-1)$$

$$(d * e) \text{ mod } \phi(n) = 1$$

Decryption

$d = \text{Private key}$

$$P = C^d \text{ mod } n$$

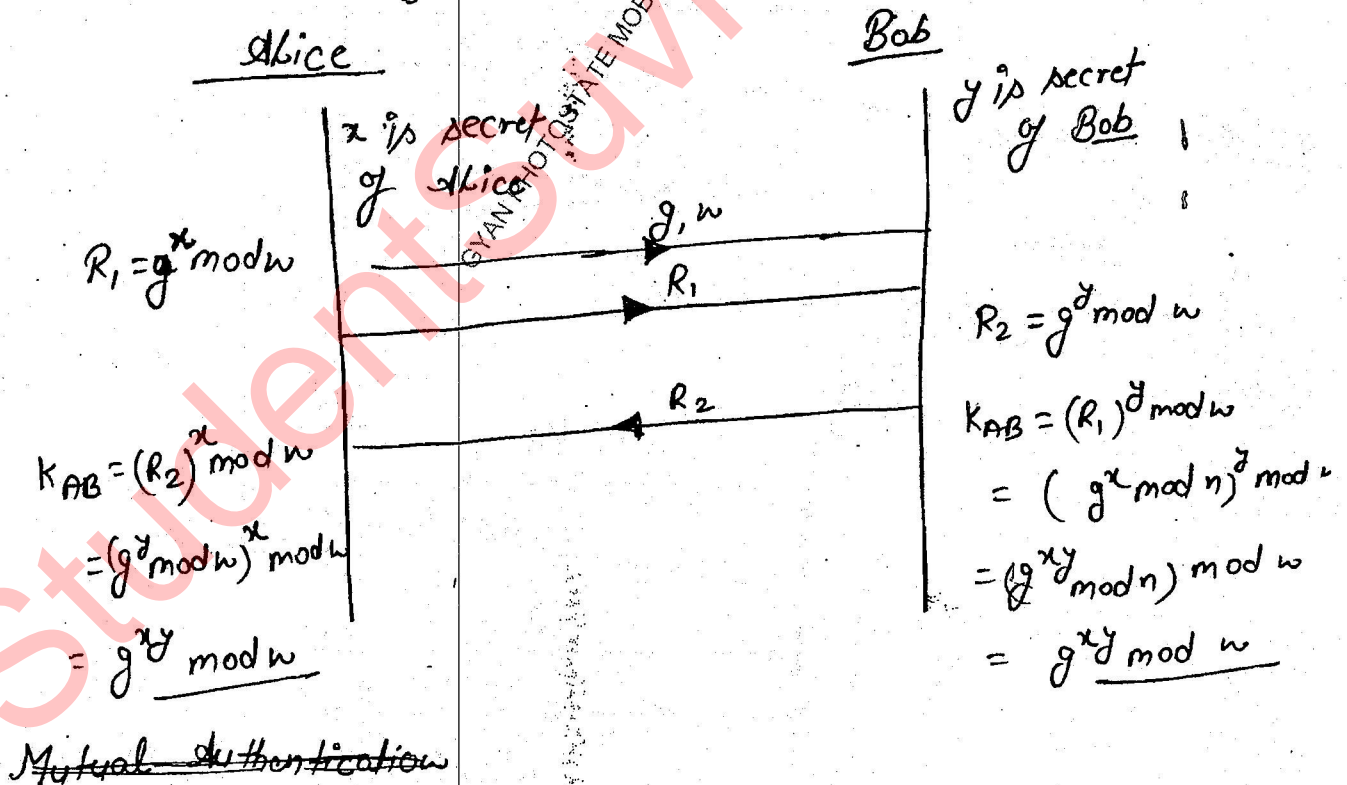
Confidentiality.

* In RSA algorithm, if sender is encrypted with Receiver's public key & Receiver is decrypted with his own private key then it is used to provide Confidentiality.

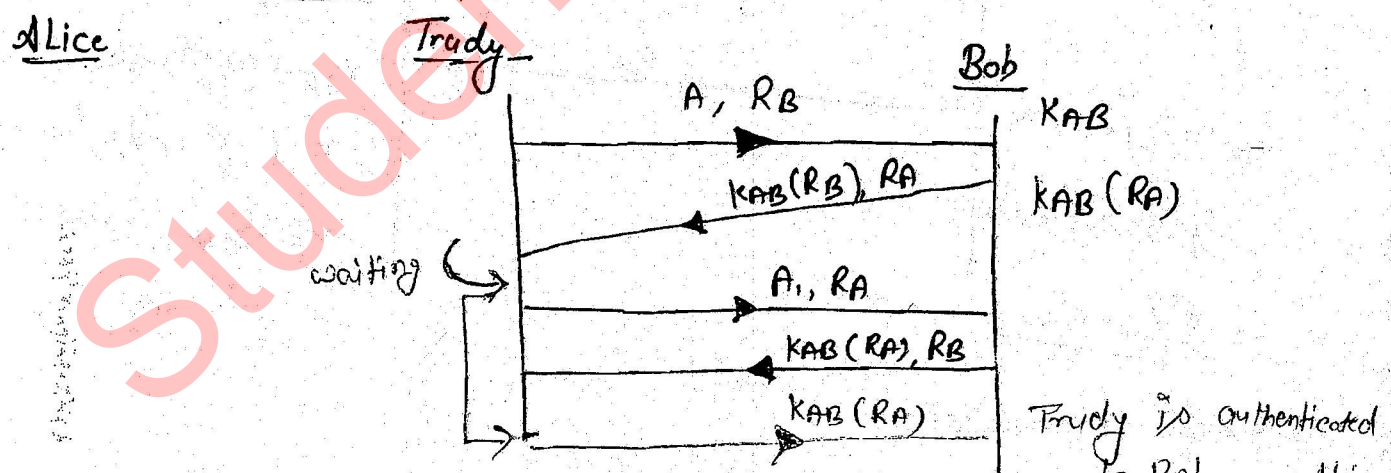
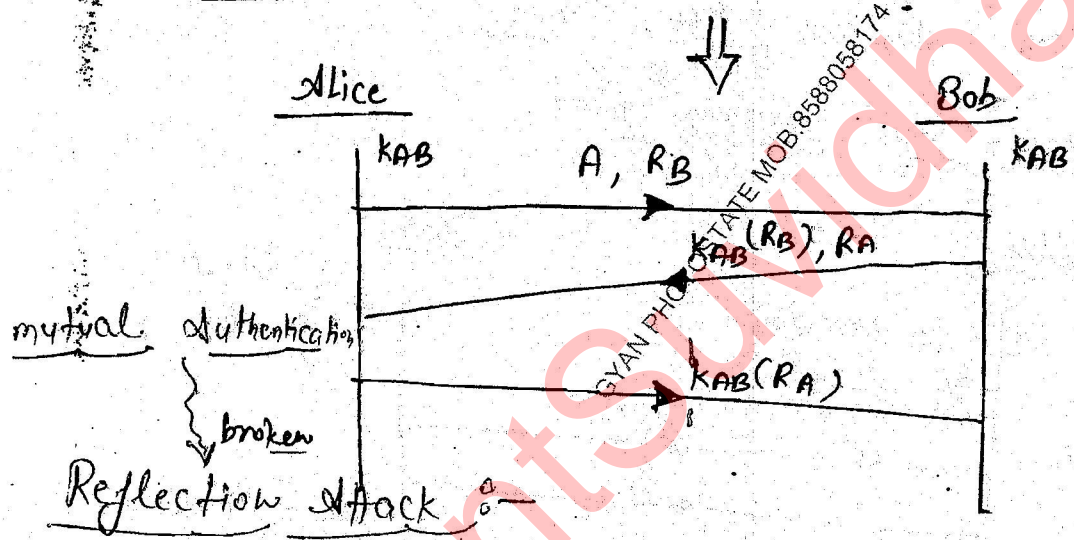
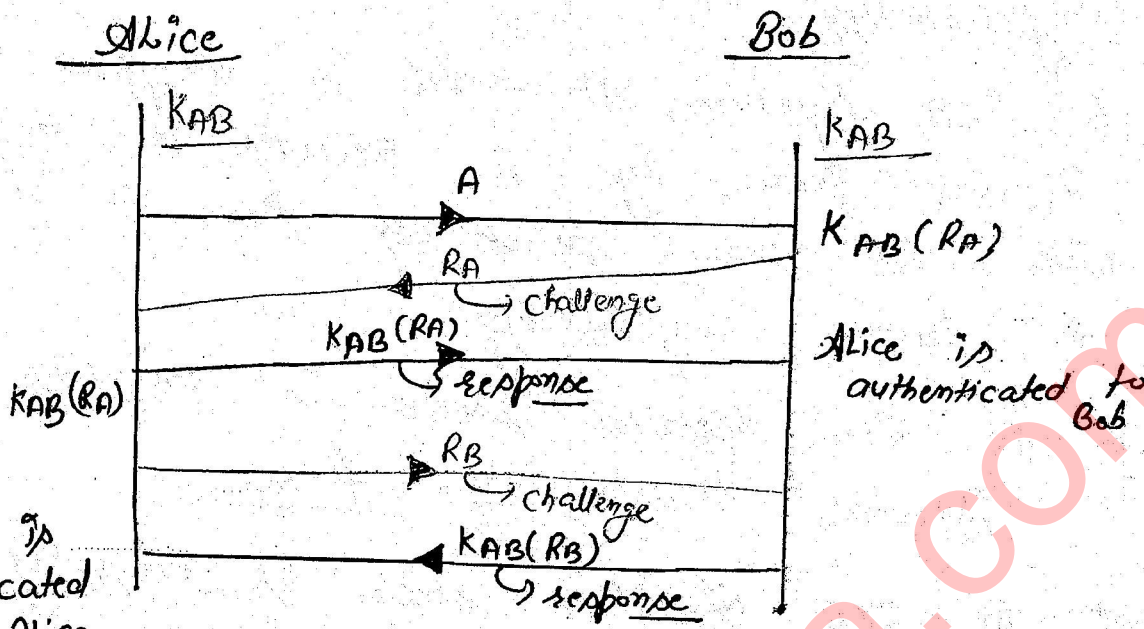
Que:- In RSA algorithm sender is encrypting with Receiver's Private key. false

Que:- In RSA algorithm, sender is encrypting with his own private key. True

Diffie-Hellman key exchange

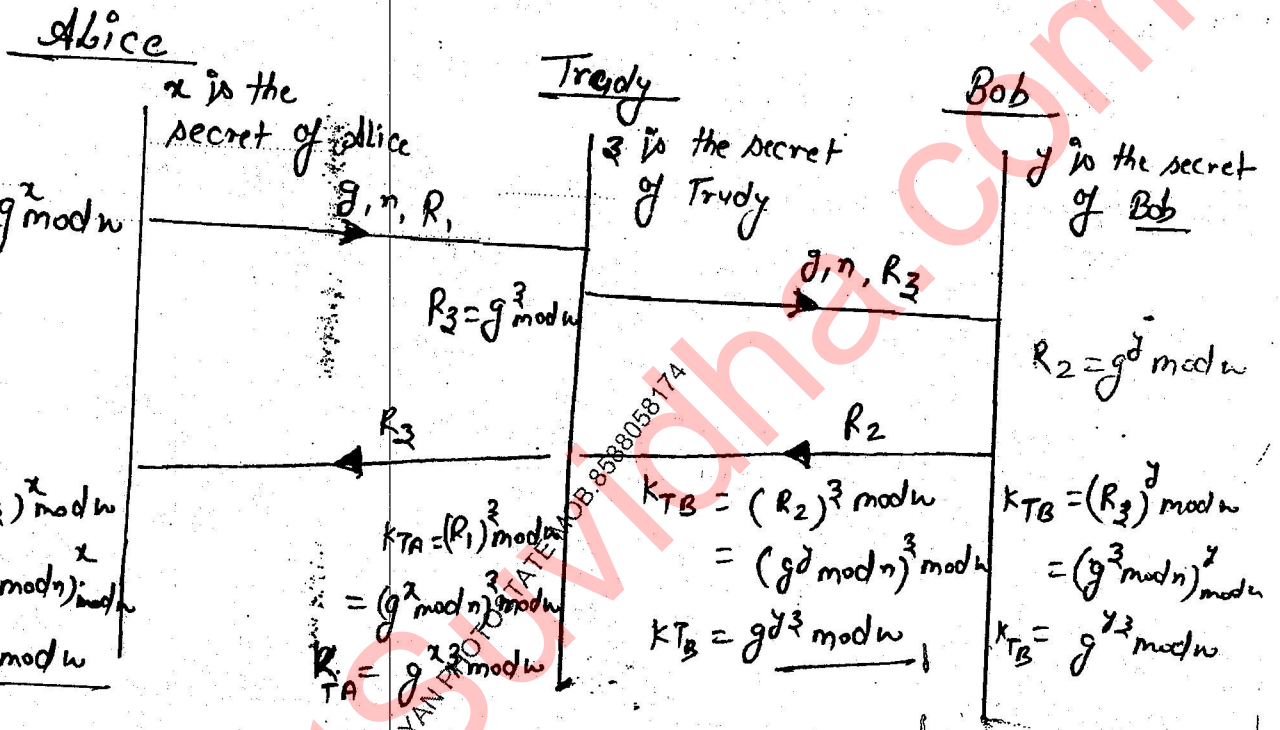


Mutual Authentication using Diffie-Hellman key exchange:



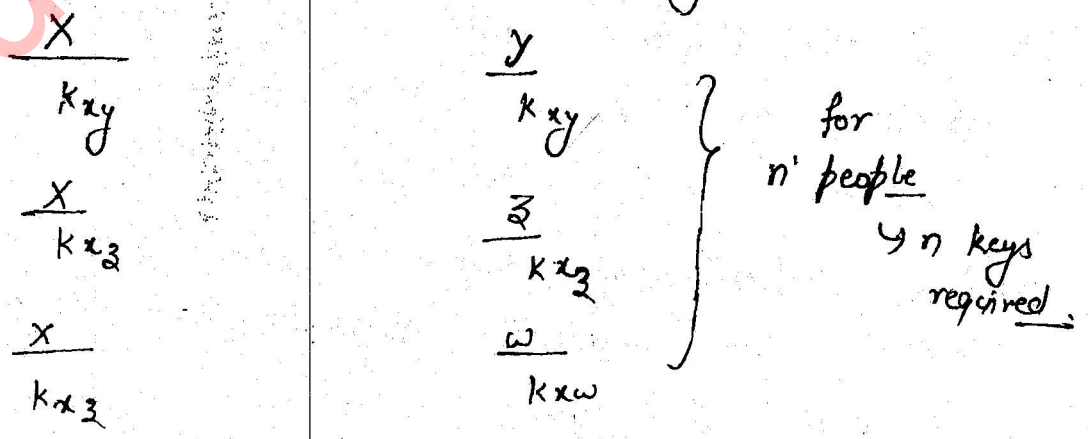
* Mutual Authentication can be broken using Reflection attack.

* whenever the packets are transmitted, if Hacker has modified the source IP then the request will go to the destination but the reply will come to Hacker.



* Diffie Hellmann key itself can be broken using MAN in middle attack.

* Drawback of Diffie Hellmann Exchange :-

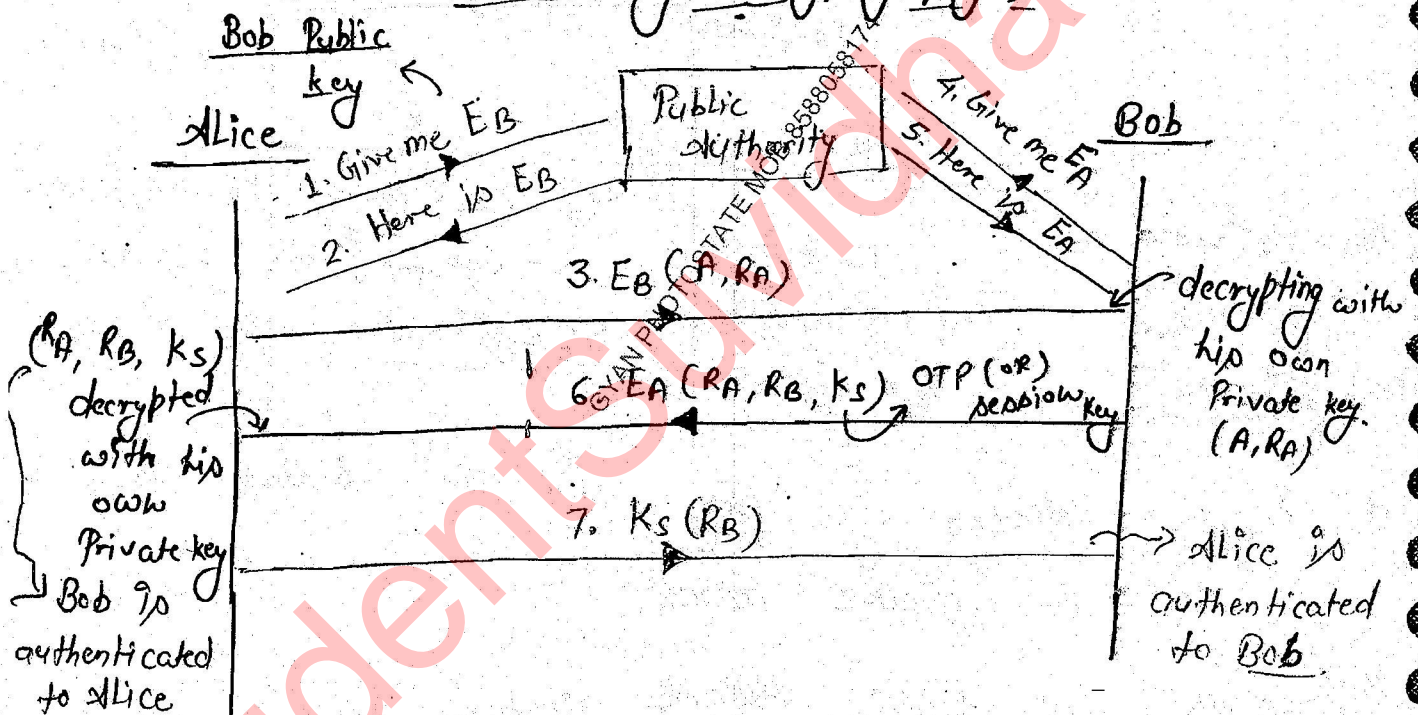


* The drawback of Diffie Hellman key is, to communicate with n people then n keys are required.

* Remembering all the key values or maintaining a database of all the key values is difficult.

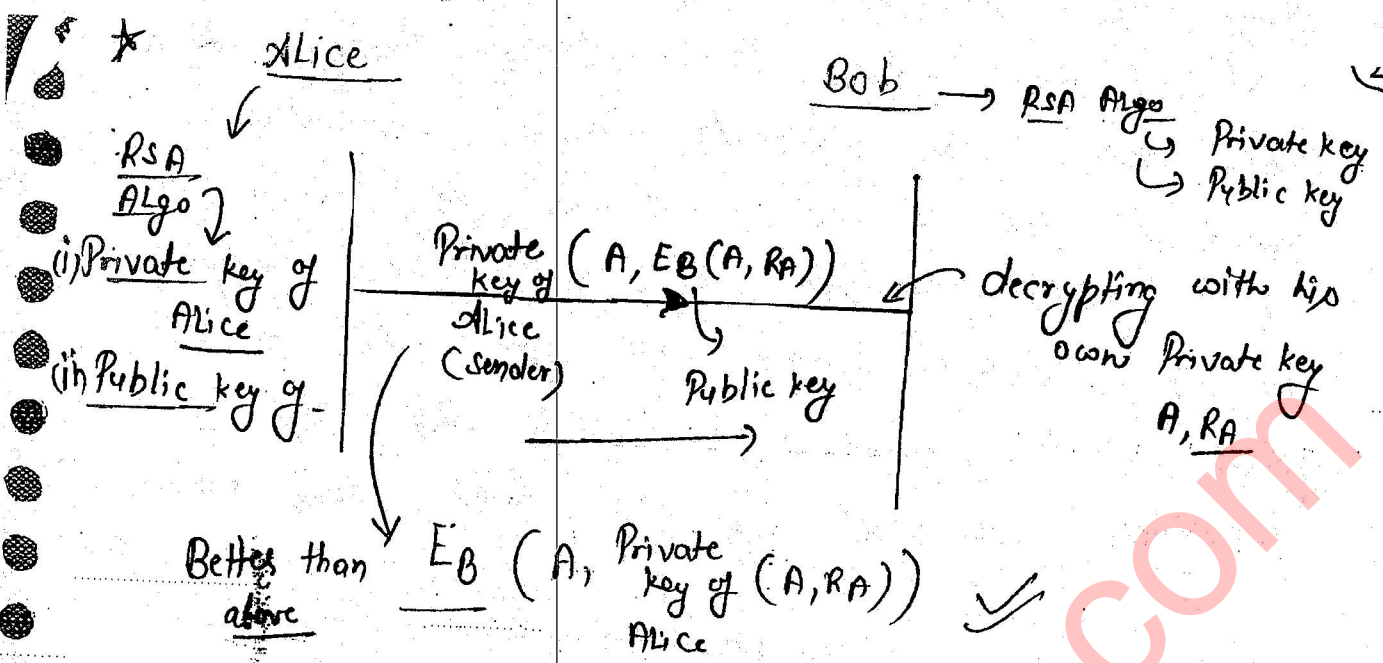
* So, the generation of the key is taken care by third party or certified authority.

⇒ Mutual Authentication using RSA algorithm :-
or (Public key cryptography)



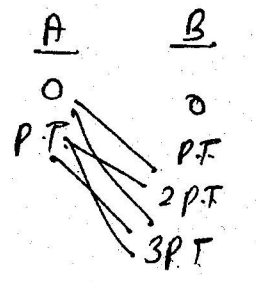
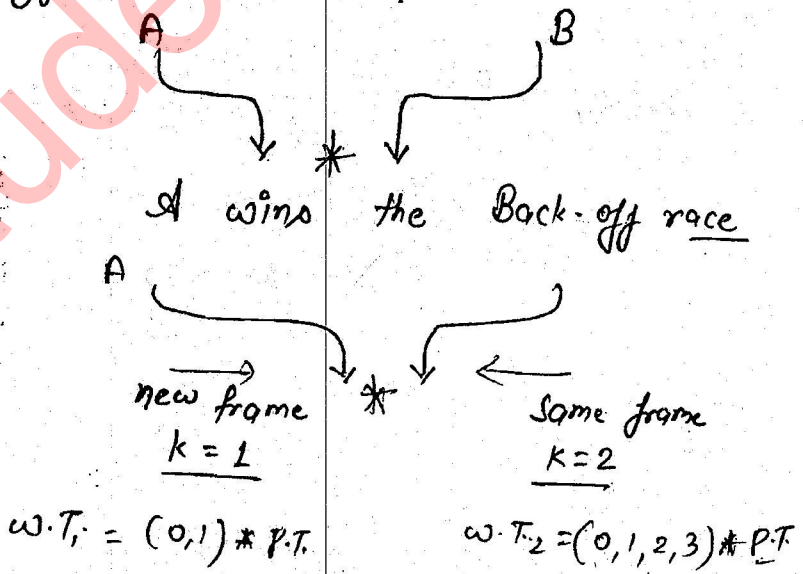
* Mutual authentication using RSA algorithm is better than mutual authentication using Diffie Hellman in terms of security.

* Mutual authentication using Diffie Hellman key is better than mutual authentication using RSA in terms of time.



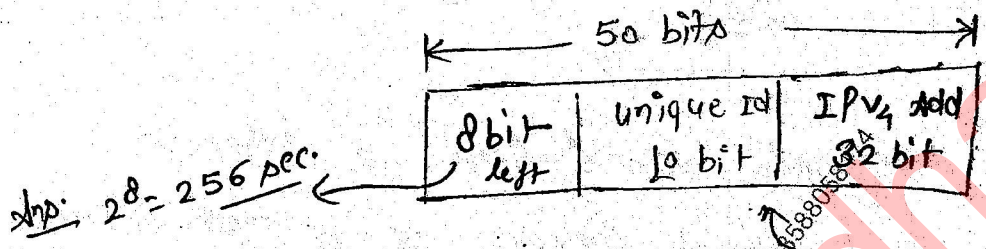
Que. **

① A and B are only two stations on Ethernet. Each has a steady queue of frames to send. Both A and B attempt to transmit a frame, collided and A wins the 1st back-off race. At the end of successful transmission by A both A and B attempt to transmit and collide. The probability that A wins the second back-off race is : ?



$$\Rightarrow \frac{3+2}{8} = \frac{5}{8}$$

② Every Host in IPv4 network, has a real time clock with battery backup. Each host needs to generate upto 1000 unique identifiers per second. Assume that each host has global unique IPv4 address. Design a 50 bit global unique ID for this purpose. After what period (in sec.) will the identifiers generated by a host wrap around



1000 unique identifiers = $2^{10} \leftarrow 10 \text{ bits}$
 $2^9 = 512$ not sufficient

Note :-

For Stop & wait Protocol L.U. = $1 * \frac{T.T.}{T.T. + 2 * P.T.}$

For GBN Protocol L.U. = $\frac{N * T.T.}{T.T. + 2 * P.T.}$

No. of frames \rightarrow window size

Ques:- Consider a store & forward packet switched network. The B.W. of each link is 10^6 bits/sec. User on host A sends a file of size 10^3 Bytes to host B. in three different ways:-

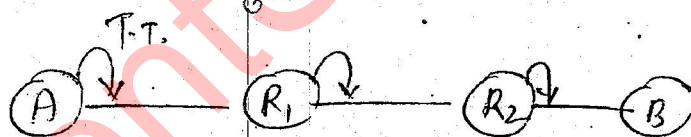
- Case:-
- (i) A single packet containing the complete file is transmitted from A to B.
 - (ii) The file is split into 10 equal parts & these packets are transmitted from A to B.
 - (iii) the file is split into 20 equal parts. (Each packet contains 100 Bytes of header along with the user data.) (Consider only transmission time & ignore processing, propagation delay)

Solⁿ:-

file = 10^3 bytes

B.W. = 10^6 bits/sec.

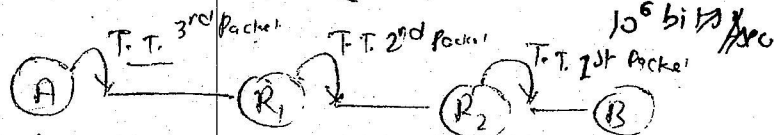
$$T.T. \text{ of Packet} = \frac{\text{Data size}}{\text{B.W.}} = \frac{(1000 + 100) * 8}{10^6 \text{ bits/sec.}} = \frac{8800}{10^6} = 8.8 \text{ millise}$$



(i) A to B = $3 * 8.8 \text{ millise} = 26.4 \text{ millise}$

(ii) $\frac{1000 \text{ bytes}}{10} = 100 \text{ bytes}$ → Each packet size
 $100 \text{ bytes header} = 200 \text{ bytes}$

$$T.T. \text{ Packet} = (100 + 100) = \frac{200 * 8}{10^6 \text{ bits/sec.}} = 1.6 \text{ millise}$$



Total time to reach B of TT₁ = $3 * T.T. = 3 * 1.6 = 4.8 \text{ millise}$
 2nd Packet = $4.8 + 1 * 1.6 =$

$$20^{\text{th}} \text{ Packet} = 4.8 + 9 * 1.6$$

$$= 19.2 \text{ millisee.}$$

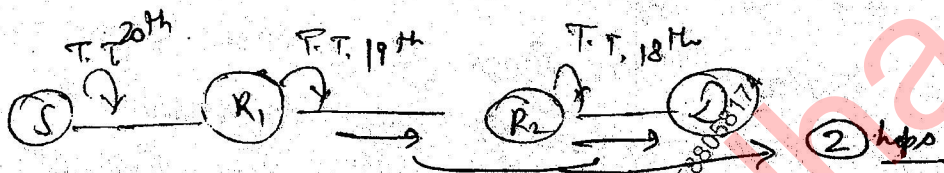
(iii)

$$\frac{1000 \text{ Bytes}}{20} = 50 \text{ Bytes} + 100 \text{ Byte} = 150 \text{ Bytes}$$

Header

^ Each Packet size

$$\text{T.T. of Each Packet} = \frac{150 * 8}{10^6 \text{ bit/Dec.}} = 1.2 \text{ milli sec.}$$



$$20 \text{ Packets} = 20 * 2 = 24 \text{ milli sec.}$$

$$19^{\text{th}} \text{ Packet} \Rightarrow 24 + 2 * 1.2 = 26.4$$